

Page (Line)	Error	Correct it to
Chapter 1:		
3 (3)	to provide	provide
13 (3)	in fact resulted ...	resulted ...
13 (26)	n_h	n_p
29 (20)	a packet	data
Chapter 2:		
22 (17,19)	Layer 6, Layer 7	Layer 5, Layer 6
25 (14)	65,534	65,535
28 (13)	205.101.0.0/19	205.101.0.0/20
Chapter 3:		
45 (11)	34	31
46 (11)	240	170
47 (9-10)	... available the transmission line but the multiplexer m has input, where available, where ...
47-48-49	(eight cases) $1 \leq j \leq m, 1 \leq j \leq n, \sum_{i=1}$	$0 \leq j \leq m, 0 \leq j \leq n, \sum_{i=0}$
48 (12)	three	five
48 (16)	$P_{n=5} = \frac{\binom{12}{3} \left(\frac{2}{3}\right)^5}{\sum_{i=0}^5 \binom{12}{i} \left(\frac{2}{3}\right)^i} = 0.44.$	$P_{n=5} = \frac{\binom{12}{5} \left(\frac{2}{3}\right)^5}{\sum_{i=0}^5 \binom{12}{i} \left(\frac{2}{3}\right)^i} = 0.72.$
51 (Fig. 3.6)	Original Bits: 11011101101011010111010101101 NRZ-Inverted:	Original Bits: 11011101101011011111010101101 NRZ-Inverted:
51 (17)	... all ... 1s and all 1s as Os a sequence of ... a weak 1 or even a 0 ...
52 (3)	1ns	1s
53 (11-12)	... difference in the sence that three different in the sense that there ...
67 (31)	4,000 KHz	4,000 Hz
68 (4)	170	160
69 (8)	Manchester NRZ	Manchester
Chapter 4:		
90 (24)	lke	like
94 (Fig. 4.9)		
94 (5)	... with the bits shift left in the register.	... with the bits in the register.
94 (14)	$x^4 + x + 1$	$x^4 + x$
100(2)	show the value of CRC	show 0000
100(Prob.9)	Data frames with ...	A total of 1000 data frames with ...
Chapter 5:		
102 (19)	... layer 2 devices called layer 1 or 2 devices such as ...
110(17)	... a user	... a user having a frame to transmit during a fair share of time, and
111 (16)	$P_X(k) = \frac{\lambda^k e^{-\lambda}}{k!}$	$P_X(k) = \frac{(\lambda t)^k e^{-\lambda t}}{k!}$

118 (11)	user 5	user 3
119 (4)	LAN 2, B3, and LAN 3	LAN 5, B4, and LAN 4
120 (1)	B2 and B3	B1 and B2
121 (Table 5.1)	Routing table for two bridges (B1 and B3) ...	Routing table for two bridges (B1 and B2) ...
121 (Table 5.1)	LAN 1 ... LAN 1 ... LAN 2 ... LAN2 LAN 1 ... LAN 1 ... LAN 2 ... LAN2	LAN 1 ... LAN 1 ... LAN 2 ... LAN2 LAN 2 ... LAN 2 ... LAN 4 ... LAN4
121 (Table 5.1)	User 4, LAN 2, User4, -, User 4, -, User 4, -	User 4, LAN 2, User4, -, User 4, -, User 4, LAN 2
121 (Table 5.2)	00-40-33-25-85-BC LAN 1	00-40-33-25-85-BC LAN 2
127 (Problem 3)	100 m	1000 m (also combine parts a and b)
127 (6)	22 K	22 KB
Chapter 6:		
147 (4)	$m = 3$	$m = 2$
156 (32)	iss	is
Chapter 7:		
178 (17)	Dijkstra	Dijkstra
178 (21)	... obvious ... the least- obvious that ... the cost of least-
180 (Table 7.2)	AED(4)	ACED(4)
180 (Table 7.2)	ACEG(6)	ACEDG(6)
208 (Prob.6)	the two servers	R1 and R7
209 (3 & 10)	R_7	R_4
Chapter 8:		
210 (27)	filed	field
211 (Fig. 8.2)	RSI	RST
218 (13)	diviationsare	diviations are
224 (17, 19, 24)	K	KB
Chapter 9:		
233 (4)	<i>terminal</i>	<i>telecommunication</i>
234 (4)	... option that option is that ...
Chapter 10:		
258 (Fig. 10.3)	$f \dots f$	$F() \dots F()$
258 (8)	52	56
258 (12)	left and right	right and left
262 (1)	1,000	1,001
271 (15)	k_4	k_5
Chapter 11:		
279 (Eq. 11.11)	$\sum_{j=0}^{\infty} (\rho_1 \rho_2 \dots \rho_i)$	$\sum_{j=0}^{\infty} (\rho_1 \rho_2 \dots \rho_j)$
287 (19)	arrival	service
289 (Eq. 11.29)	p	p_0
291 (<i>Solution</i>)	... 0.095 ... 0.38 ... 3.38 0.12 ... 0.48 ... 3.48 ...
300 (Fig. 11.18)	$\mu_m \dots \mu_{m-1}$	$\mu_{m-1} \dots \mu_m$
303 (22)	$P_i \lambda_i / \mu_i$	$P_i \lambda / \mu_i$
312 (Prob. 12)	$\mu_1 = 4, \mu_2 = 3, \mu_3 = 5, \text{ and } \mu_0 = 2$	$\mu_1 = 4, \text{ and } \mu_2 = 3$ (so, eliminate units 3 and 4 from Figure 11.26, and change "four" in Part (a) to "two")
313 (5)	μ_0	μ_1
Chapter 12:		
316 (6)	Traffic regulation traffic ...	Traffic regulation ...

319 (12)	$1/g$	i/g
320 (26)	a grant has arrived	i grants have already been assigned to i packets in the main buffer
327 (17)	$1/\mu_i$	μ_i
330 (Eq. 12.16)	$\dots + \left(\frac{1}{\mu} + \theta_i\right)$	$\dots + \theta_i$
330 (9)	θ_i	θ_i
331 (Eq. 12.19)	$\dots + \frac{1}{\mu} + \dots$	$\dots + \dots$
333 (29)	20	30
334 (Fig. 12.13)	WFQ: A-1 B-1 D-1 B-2 B-3 C-1 C-2 C-3 D-2 D-3 A-2 A-3	WFQ: A-1 B-1 D-1 D-2 A-2 B-2 C-1 C-2 C-3 D-3 A-3 B-3
336 (35)	PHHB	PHB
345 (4)	Assume	Consider only non-queueing states with the assumption of
345 (6)	\dots 0,1,2, or 3 grants are allocated to arriving packets	\dots 0, or 1 grants are left in the grant buffer
345 (Problem 3)	(Do the same corrections as mentioned above in Problem 2)	
345 (31), 346 (3)	$\dots 1/\mu = 1/\text{ms} \dots r = 0.5/\text{ms} \dots$	$\dots 1/\mu = 1 \text{ ms} \dots r = 0.5 \text{ ms} \dots$
348 (6)	20	30
Chapter 13:		
355 (15)	Delta	Omega
356 (4)	network self-routing	network is self-routing
Chapter 14:		
390 (Fig. 14.8)	$\lambda_{2,4}$	$\lambda_{1,4}$
394 (Eq. 14.4)	$P_b = 1 - (1 - p^r)^{\lambda_n}$	$P_b = [1 - (1 - p)r]^{\lambda_n}$
Chapter 15:		
403 (4)	notingh	noting
406 (8)	are first formed	are formed
419 (6)	when	before the
420 (4,21,22)	d^{k-j}	d^{r-j}
420 (18)	$2 \log_d n$	$\log_d n$
421 (19 & 23)	$A_{j+1} = A_k \dots A_{k-(j-1)} 01 \dots 1$ $a_{j+1} = a_k \dots a_{k-(j-1)} 10 \dots 0$	$A(j+1) = A_k \dots A_{k-(j-1)} 01 \dots 1$ $a(j+1) = a_k \dots a_{k-(j-1)} 10 \dots 0$
422 (13)	$\log_{16} 2$	$\log_2 16$
427 (1)	protocols is MOSPF,	protocols, MOSPF is
429 (9)	IPS 2	ISP 2
Chapter 16:		
436 (5)	134.43..0.1 \dots 134,43.0.10	134.43.0.1 \dots 134.43.0.10
Chapter 17:		
453 (22)	$D_b \leq D$	$D_b \geq D$
455 (7)	also $\frac{\partial D}{\partial x_1} = 0, \dots, \frac{\partial D}{\partial x_N} = 0$. The result of solving the above $N + 2$ equations can be summarized as follows. For N =even:	given the following arrangements when N is odd or even:
455 (10)	$\hat{x}_i = \hat{x}_{N+1-i}$	$\hat{x}_i = -\hat{x}_{N+1-i}$
455(Eq 17.12-13)	(Eliminate Equations 17.12 and 17.13 and replace them with:)	the optimal values, Δ_{opt} and D_{opt} , can be obtained as results of solving the two differential equations.
457 (Eq. 17.11)	pixels	3-color pixel bundles
458 (Eq. 17.14)	$\dots \cos\left(\frac{\pi i(2x+1)}{2N}\right) \cos\left(\frac{\pi i(2y+1)}{2N}\right)$	$\dots \mathbf{P}[x][y] \cos\left(\frac{\pi i(2x+1)}{2N}\right) \cos\left(\frac{\pi j(2y+1)}{2N}\right)$
458 (Eq. 17.14)	$C(i), C(j) = 0$ otherwise	$C(i), C(j) = 1$ otherwise
458(24) & 459(2)	Equation (17.15)	Equation (17.14)

459 (Eq. 17.15)	$\dots \cos\left(\frac{\pi i(2i+1)}{2N}\right) \cos\left(\frac{\pi i(2j+1)}{2N}\right)$	$\dots \mathbf{T}[i][j] \cos\left(\frac{\pi i(2i+1)}{2N}\right) \cos\left(\frac{\pi i(2j+1)}{2N}\right)$
460 (Fig. 17.8 b)	-60 1 0 0 0 0 0 0	-60 0 0 0 0 0 0 0
460(2) & 474(20)	decrease	increase
462 (5)	the P frame	the B frame
465 (2)	0.522	1.74
465 (3)	$1,6000 \times 0.522 = 8,352$	$1,6000 \times 1.74 = 27,840$
466 (13,17,30)	0.522	1.74
467 (6)	R_i	R
471 (Table 17.2)	(add this between rows 4 and 5)	0 0101 00000
471 (Table 17.2)	111 0100 00000 0011 0001 10000 011 1101 10010	111 0100 00101 0010 1001 10000 010 1101 10010
471 (Table 17.2)		(eliminate last row:) 1 0001 00001
473 (5)	and a convolving pulse $c(t)$ of $\tau = 1$ volt between -1 and +1	with pulse width $\tau = 0.5$
473 (Prob. 5)	12-level optimal nonuniform	16-level optimal uniform
473 (31)	$f_X(x) = \begin{cases} 0.2 & -2 \leq X < -1 \\ 0.2 & -1 \leq X < 0 \\ 0.2 & 0 \leq X < +1 \\ 0.2 & +1 \leq X < +2 \end{cases}$	$f_X(x) = \begin{cases} 0.25 & -2 \leq X < -1 \\ 0.25 & -1 \leq X < 0 \\ 0.25 & 0 \leq X < +1 \\ 0.25 & +1 \leq X < +2 \end{cases}$
474 (2)	$f_Y(y) = \begin{cases} 0.3 & -2 \leq Y < -1 \\ 0.1 & -1 \leq Y < 0 \\ 0.4 & 0 \leq Y < +1 \\ 0.2 & +1 \leq Y < +2 \end{cases}$	$f_Y(y) = \begin{cases} 0.3 & -2 \leq Y < -1 \\ 0.1 & -1 \leq Y < 0 \\ 0.4 & 0 \leq Y < +1 \\ 0.2 & +1 \leq Y < +2 \end{cases}$
474 (14)	eight-level	four-level
474 (15)	$Q(x) = \begin{cases} +3 & 1 < x \leq 2 \\ +3 & 0 < x \leq 1 \\ +3 & -1 < x \leq 0 \\ +3 & -2 < x \leq -1 \end{cases}$	$Q(x) = \begin{cases} +1.5 & 1 < x \leq 2 \\ +0.5 & 0 < x \leq 1 \\ -0.5 & -1 < x \leq 0 \\ -1.5 & -2 < x \leq -1 \end{cases}$
Chapter 18:		
495 (Fig. 18.9)	RCTP	RTCP
502(Fig. 18.14)	Chunk n Data, Chunk 1 \dots Chunk n Data, Chunk 1	Chunk 1 Data, Chunk 1 \dots Chunk n Data, Chunk n
504 (15)	s that	is that
505 (17)	$F_X(x)$	$F_X(x)$
508 (Prob. 7)	\dots each.	\dots each. Assume each sample is encoded by 8 bits.
Appendix C:		
577 (Eq. C.19)	$\dots \sqrt{2\pi V[X]^2}$	$\dots \sqrt{2\pi V[X]}$
580 (20)	equal	equal to